



Engineering **Highway Hierarchies**

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`http://algo2.iti.uka.de/schultes/hwy/`

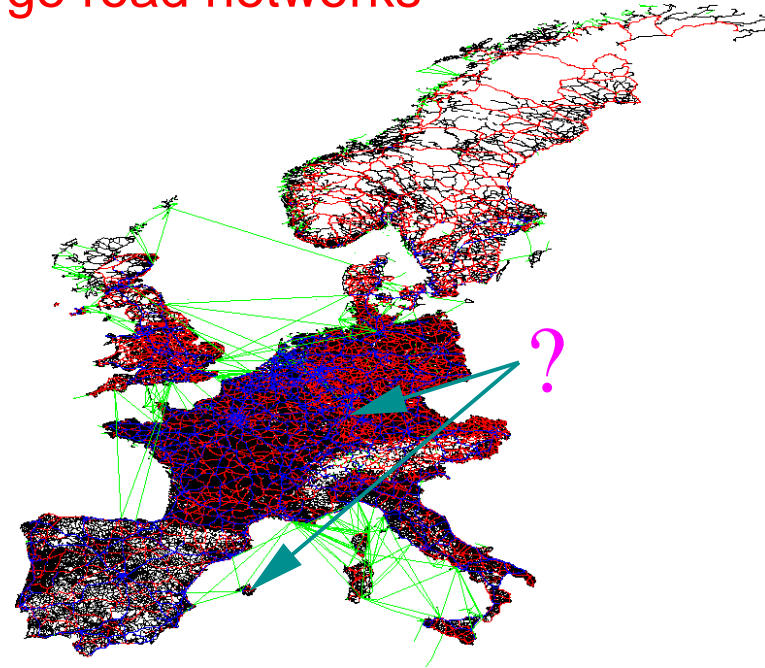
Zürich, September 11, 2006



Route Planning

Goals:

- exact shortest (i.e. fastest) paths in large road networks
- fast queries
- fast preprocessing
- low space consumption



Applications:

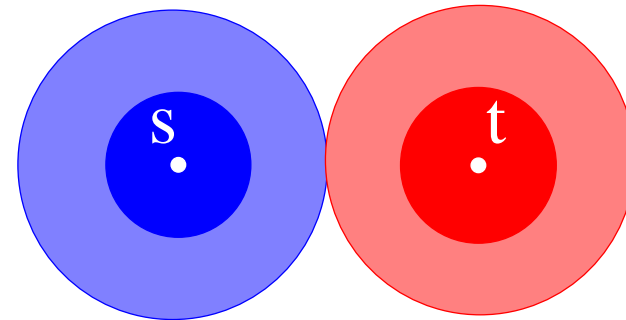
- route planning systems in the internet
- car navigation systems
- ...



Our Approach: Highway Hierarchies¹

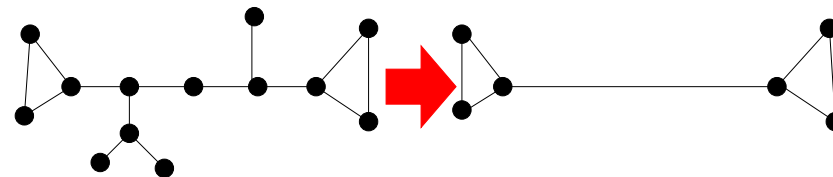
□ complete search within a local area

□ search in a (thinner) highway network

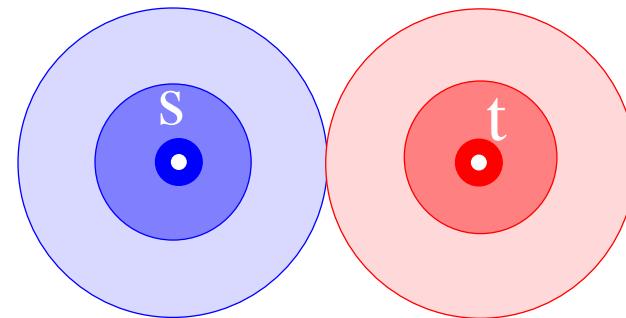


= minimal graph that preserves all shortest paths

□ contract network, e.g.,



□ iterate \rightsquigarrow highway hierarchy



¹presented at ESA 2005



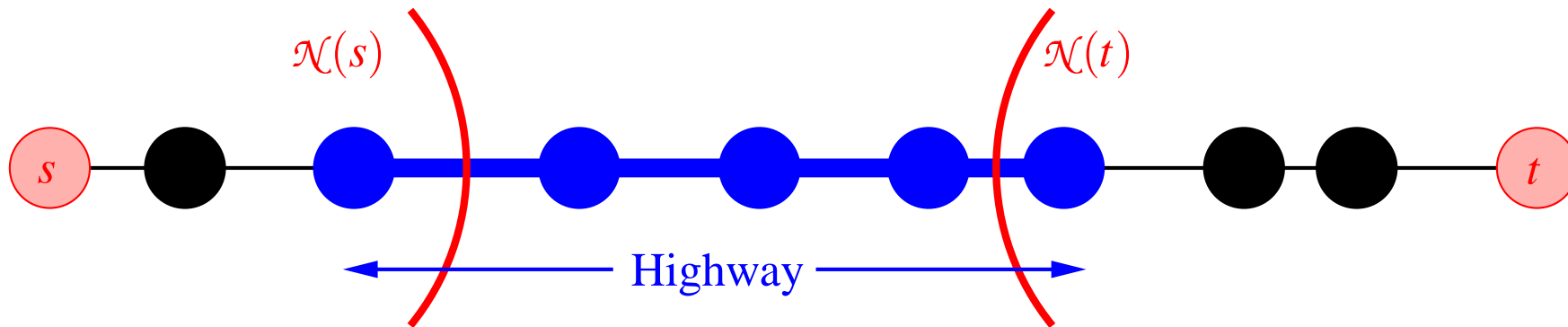
Local Area

- choose **neighbourhood radius** $r(s)$
(by a heuristic)
- define **neighbourhood** of s

$$\mathcal{N}(s) := \{v \in V \mid d(s, v) \leq r(s)\}$$



Highway Network



Edge (u, v) belongs to **highway network** iff there are nodes s and t s.t.

- (u, v) is on the “canonical” shortest path from s to t
- and
- (u, v) is not entirely within $\mathcal{N}(s)$ or $\mathcal{N}(t)$



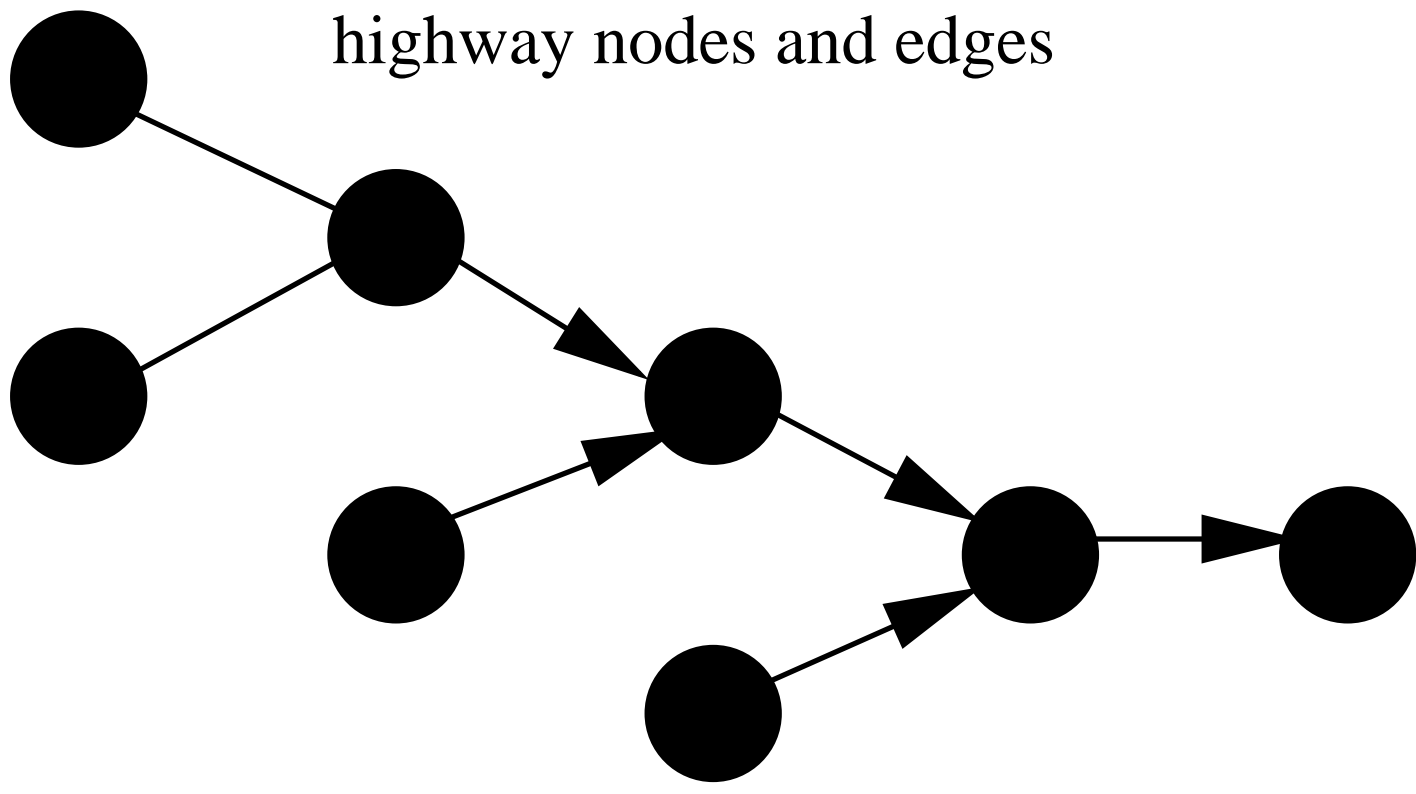
Improvements

- local area definition **more flexible**
- support** of directed graphs
- simpler**, yet **more general** and **more effective** contraction
- simpler** query algorithm
- faster** preprocessing, **faster** queries, **less** memory usage
- per-instance worst case performance **guarantees**



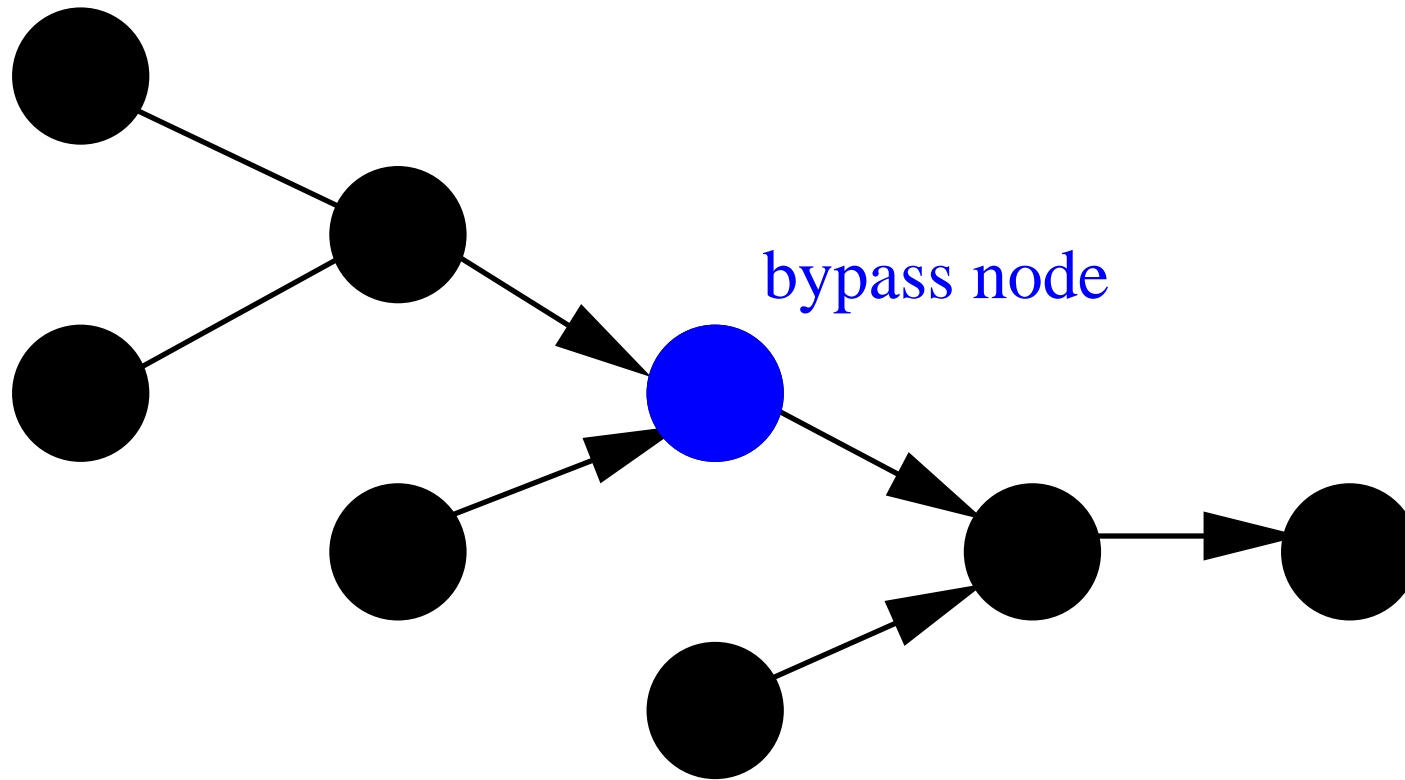
Contraction

highway nodes and edges



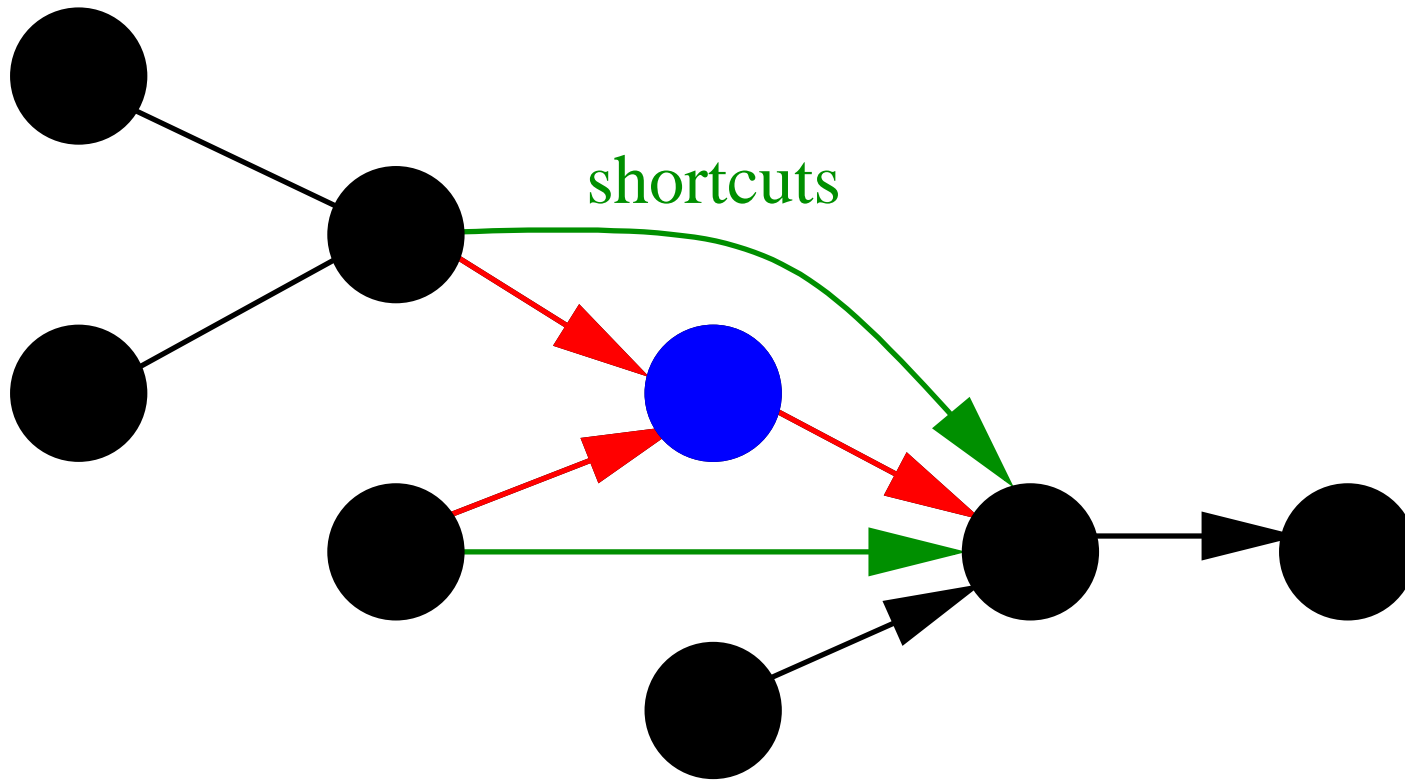


Contraction



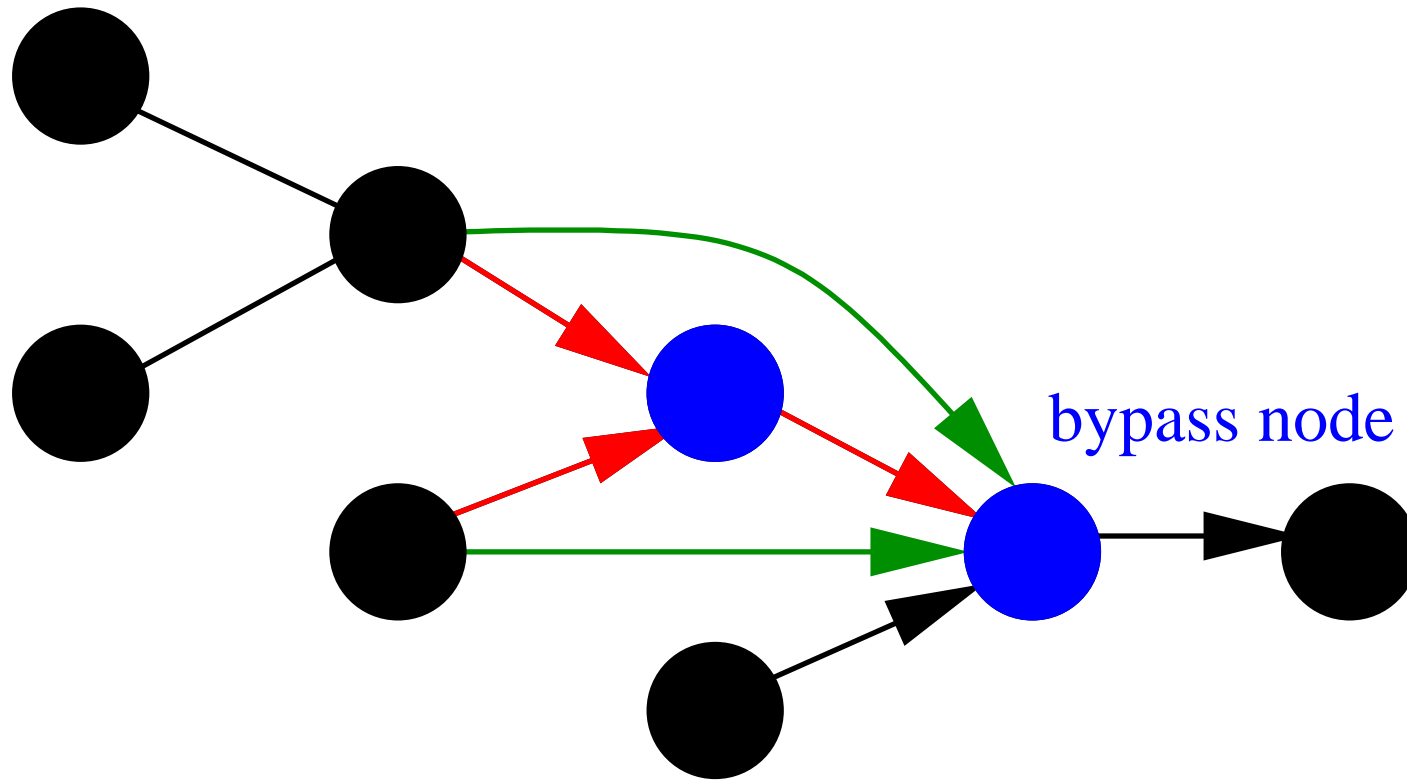


Contraction



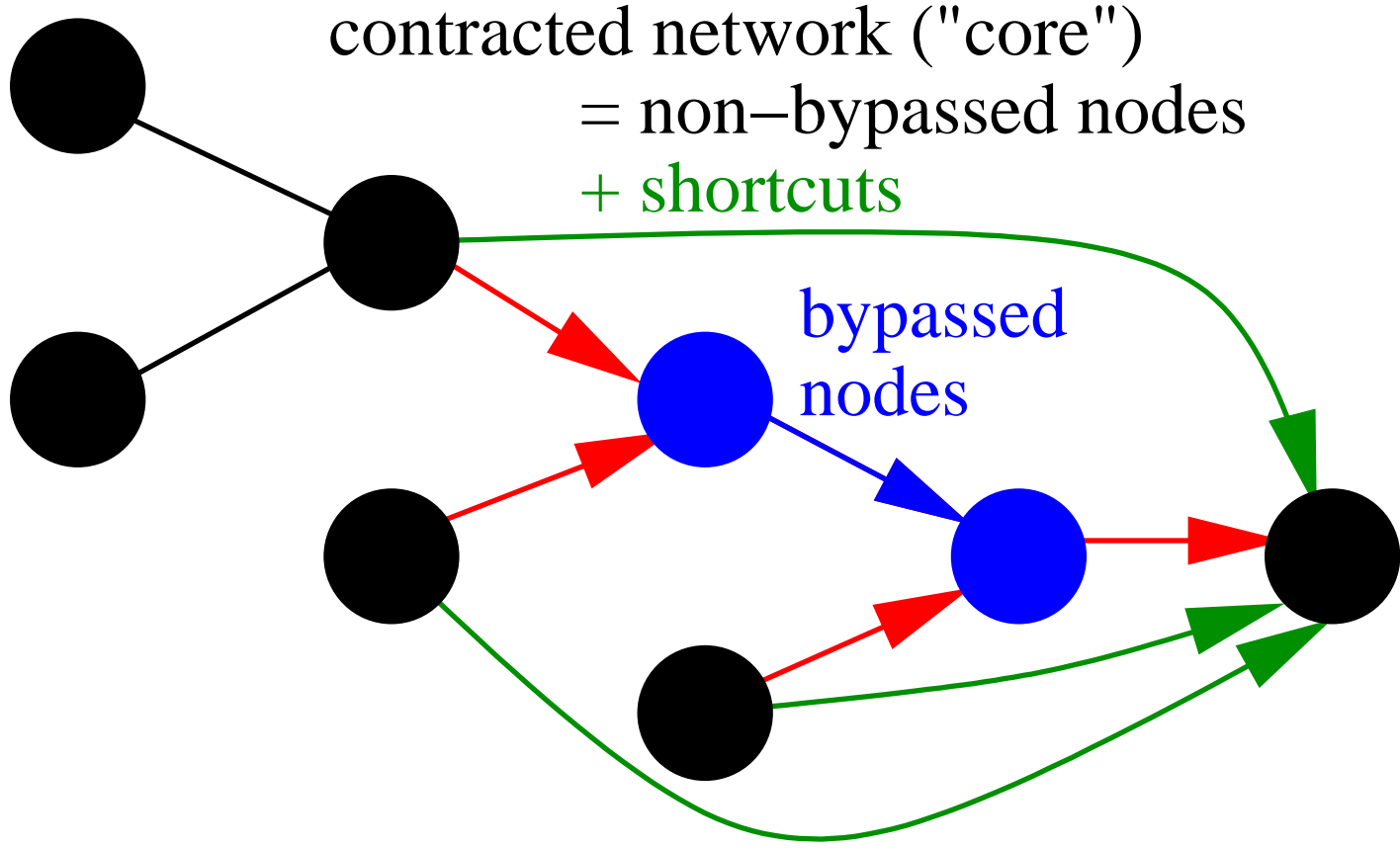


Contraction





Contraction





Contraction

Which nodes should be **bypassed**?

Use some **heuristic** taking into account

- the **number of shortcuts** that would be created and
- the **degree** of the node.



Query

Bidirectional version of **DIJKSTRA's** algorithm

+ **restrictions** on the edges that are relaxed

+ a very simple **abort** criterion

↙ search space size **increases**² by $\approx 50\%$
running time **decreases**² by $\approx 50\%$

²compared to ESA 2005

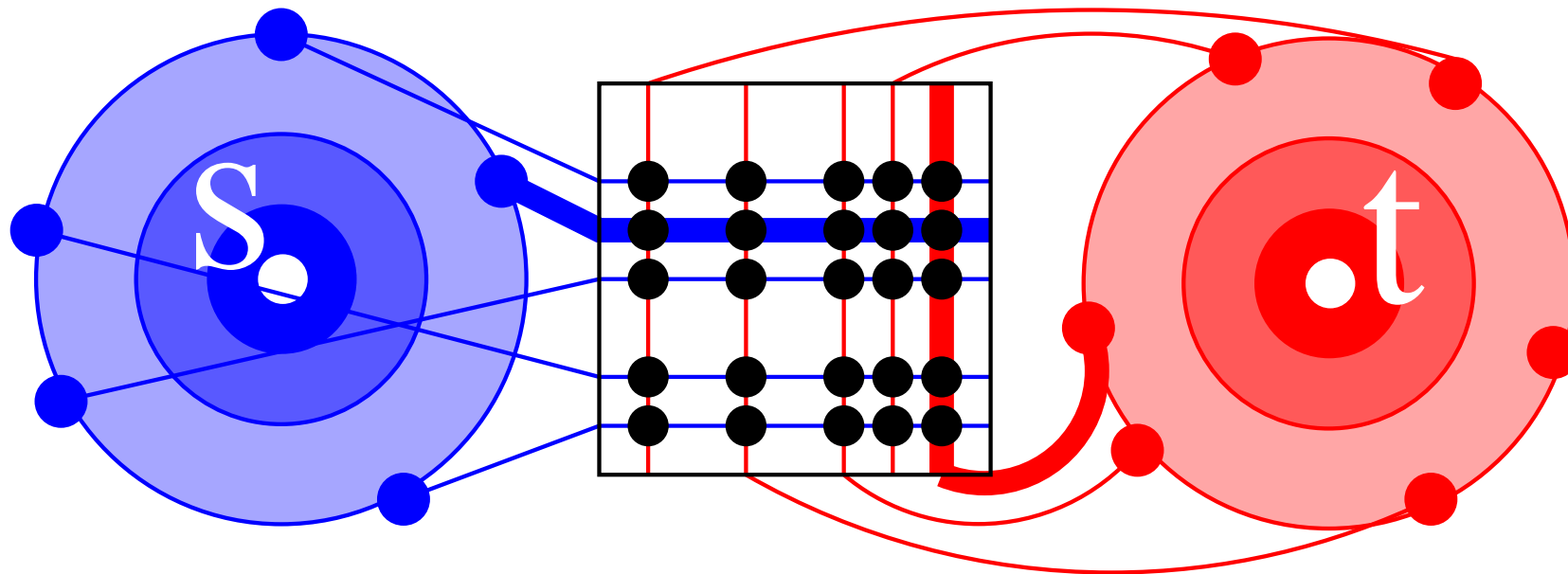


Distance Table: Construction

- Construct **fewer levels**. e.g. 4 instead of 9
- Compute an **all-pairs distance table**
for the core of the topmost level L . $13\,465 \times 13\,465$ entries



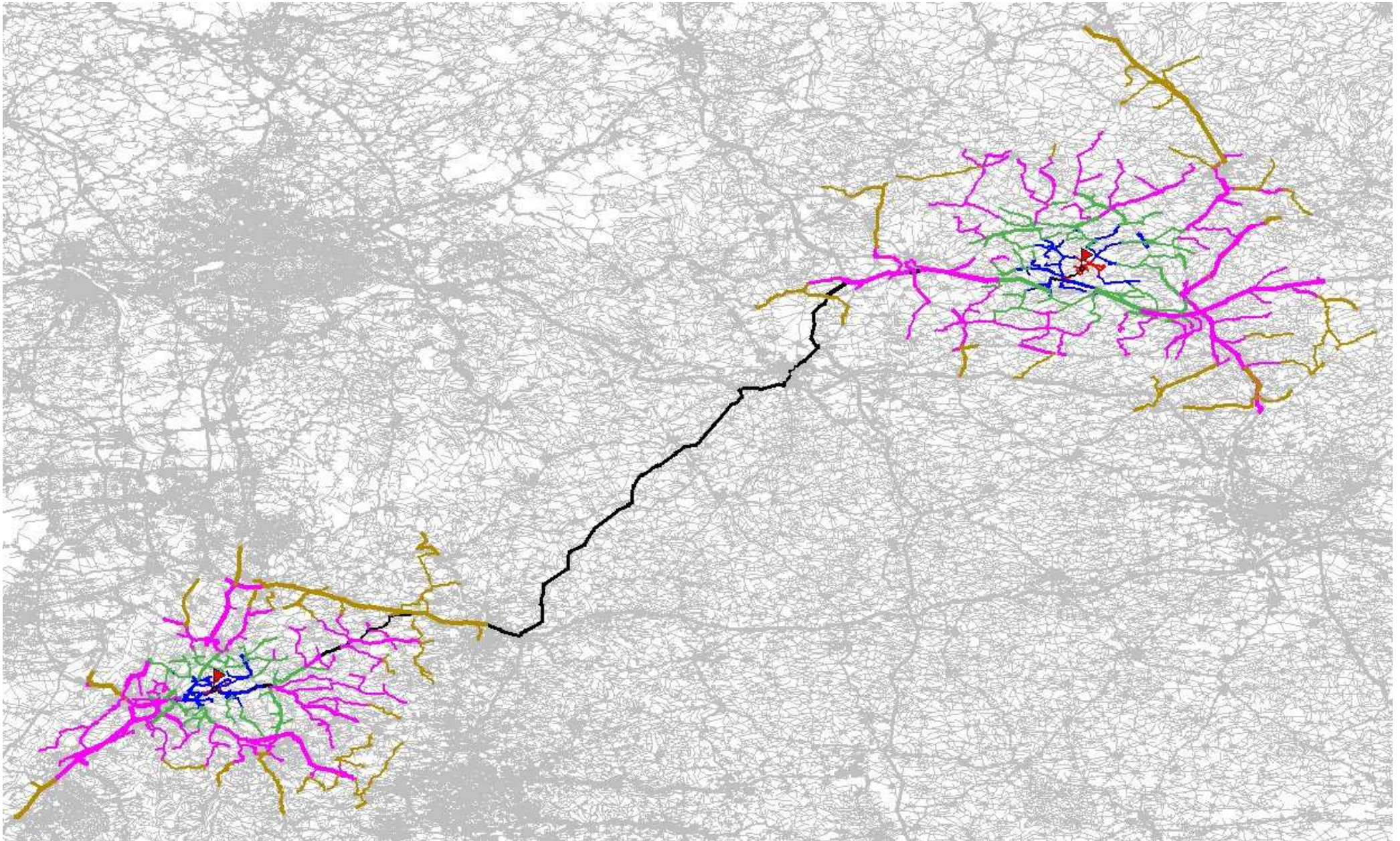
Distance Table: Query



- Abort the search when all entrance points in the core of level L have been encountered. ≈ 55 for each direction
- Use the distance table to bridge the gap. $\approx 55 \times 55$ entries



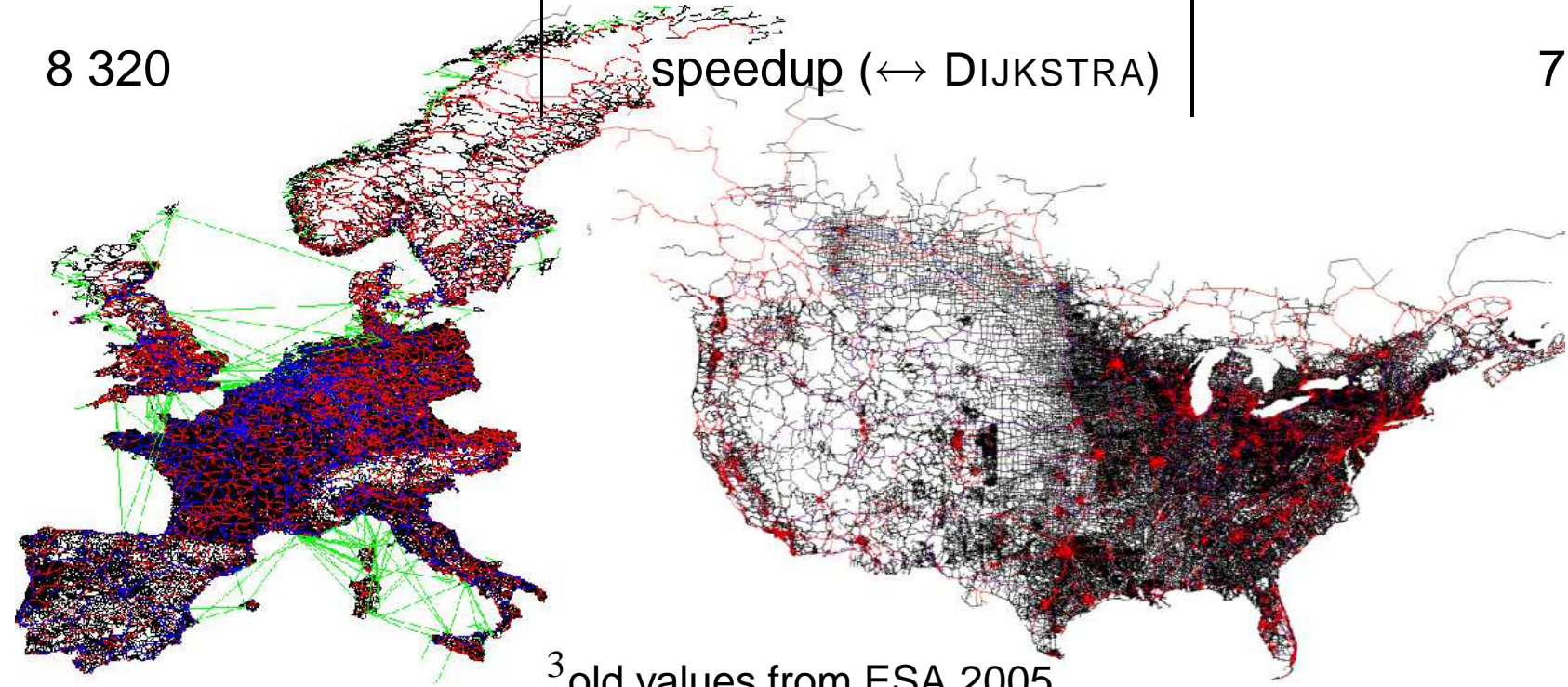
Distance Table: Search Space Example





Experiments

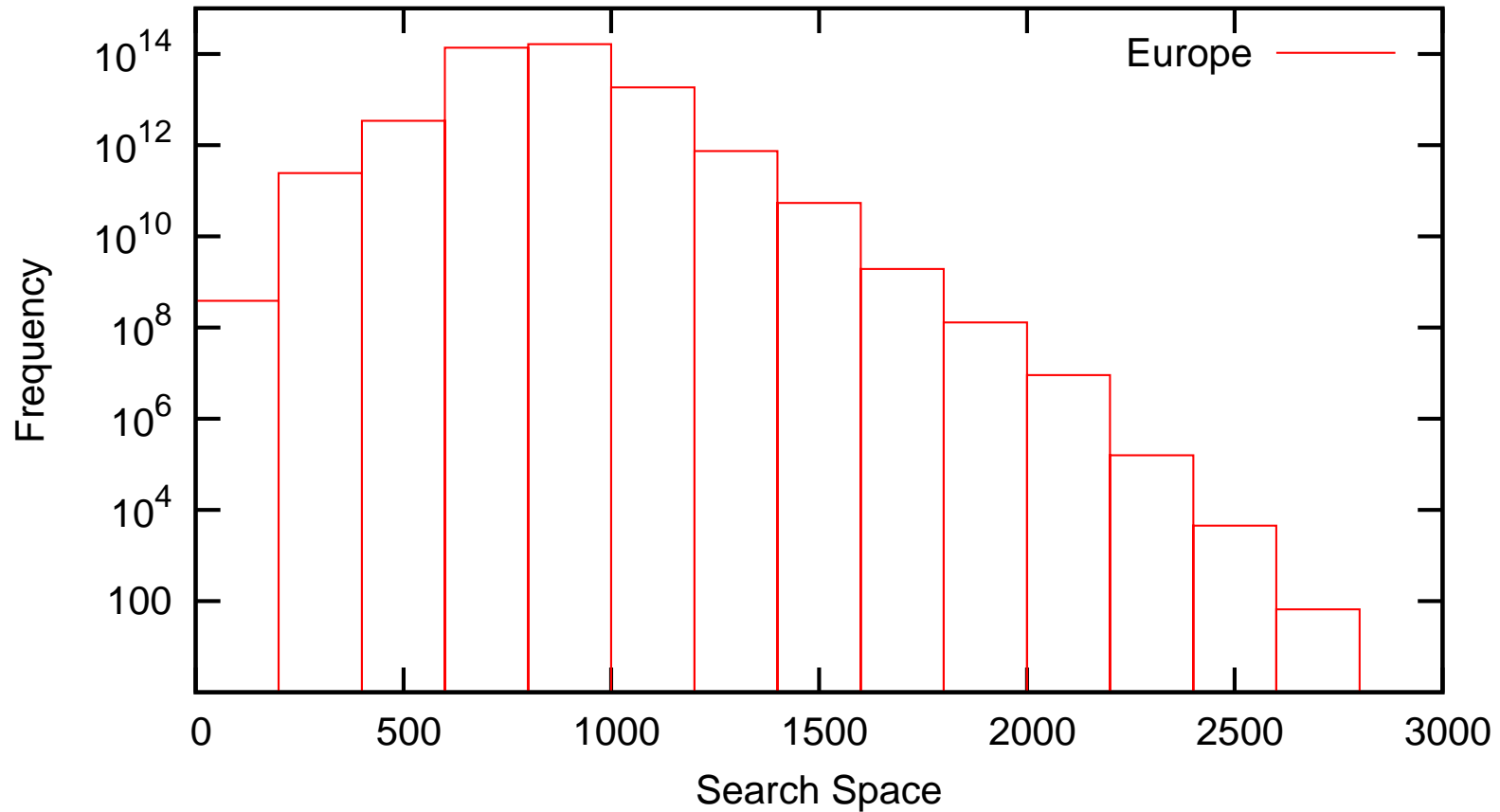
W. Europe (PTV)			USA/CAN (PTV)	
18 029 721		#nodes	18 741 705	
42 199 587		#directed edges	47 244 849	
15	[161] ³	construction [min]	20	
0.76	[7.38] ³	search time [ms]	0.90	
8 320		speedup (↔ DIJKSTRA)	7 232	



³old values from ESA 2005



Worst Case Costs⁴



Worst Case for Europe: 2 737 settled nodes (< 0.016% of all nodes)

⁴using a new feature that limits the maximum shortcut length



Summary

- exact** routes in **large** road networks (**directed!**) 18 million nodes
- fast** search 0.76 ms
 - ↪ **cheap**, **energy** efficient processors in **mobile devices**
 - ↪ low **server** load
 - ↪ lots of room for **additional functionality**
- per-instance worst case **guarantees** search space $\leq 2^{737}$
- fast** preprocessing 15 min
- low space consumption** 17–68 bytes/node



Work in Progress

- combination with a **goal directed** approach (landmarks)
joint work with [D. Delling, D. Wagner]⁵

- computation of **$M \times N$ distance tables**
(e.g. 10000×10000 table in one minute)
joint work with [S. Knopp, F. Schulz, D. Wagner]^{5,6}

- storing all entrance points into the core of the topmost level
↪ **very fast queries** ($< 20\mu s$)
joint work with [H. Bast, S. Funke, D. Matijevic]⁷

⁵Universität Karlsruhe, Algorithmik I

⁶PTV AG, Karlsruhe

⁷Max-Planck-Institut für Informatik, Saarbrücken

Future Work

fast, **local updates** on the highway network
(e.g. for traffic jams)

implementation for **mobile devices**
(flash access, ...)

multi-criteria shortest paths

joint work with [M. Müller-Hannemann, M. Schnee]⁸

...



⁸Technische Universität Darmstadt