

Algorithm Engineering for Large Graphs

Engineering Route Planning Algorithms

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Online Topological Ordering for Dense DAGs

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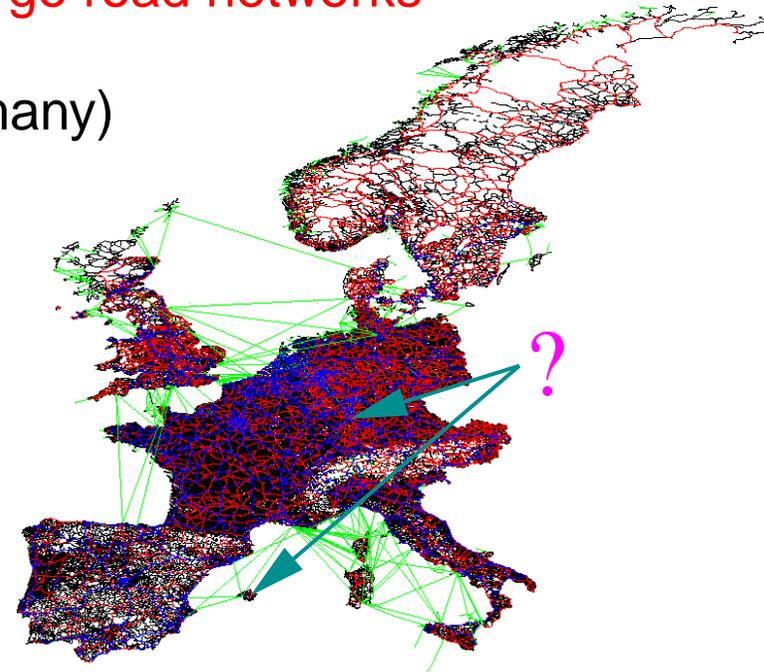
Freiburg, July 4, 2007



Route Planning

Goals:

- exact** shortest (i.e. fastest) paths in **large road networks**
- fast queries** (point-to-point, many-to-many)
- fast preprocessing**
- low space** consumption
- fast update** operations



Applications:

- route planning systems in the internet, car navigation systems,
- traffic simulation, logistics optimisation



Overview

HH Star
goal-directed
[DIMACS 06]

Transit Node Routing
very fast queries
[DIMACS 06, ALENEX 07,
Science 07]

Highway Hierarchies
foundation
[ESA 05, ESA 06]

Hwy-Node Routing
allow edge weight changes
[WEA 07]

Many-to-Many
compute distance tables
[ALENEX 07]

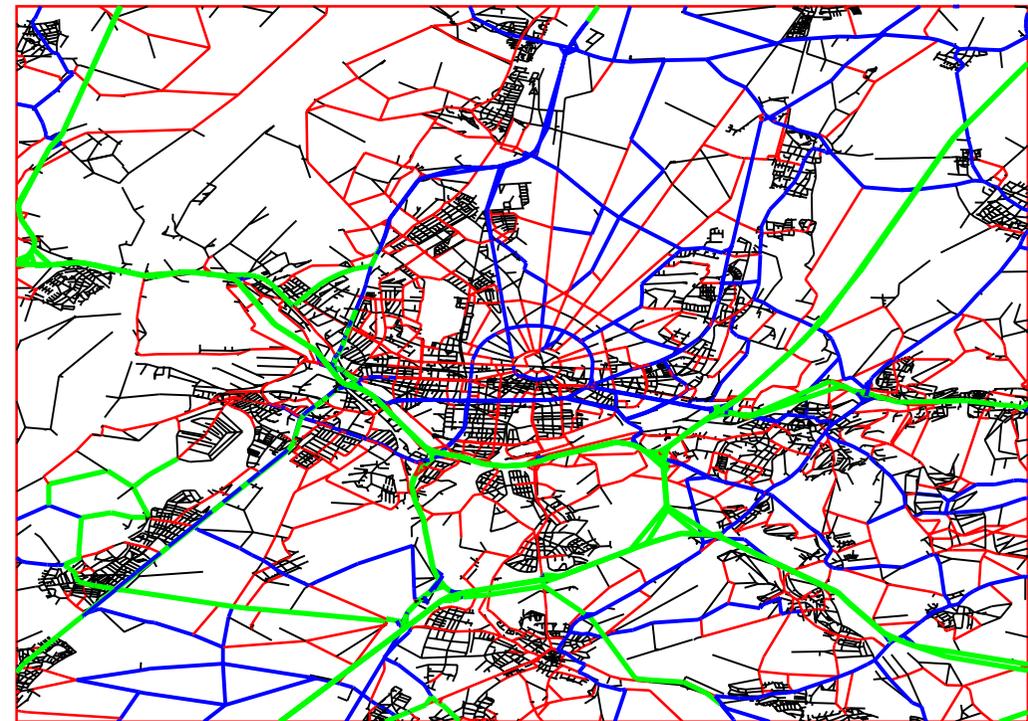


Highway Hierarchies

[ESA 05, ESA 06]

Construction: iteratively **alternate** between

- removal** of low degree **nodes**
- removal** of **edges** that only appear on shortest paths close to source or target

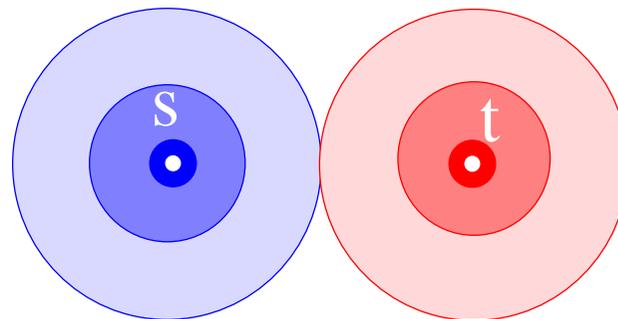
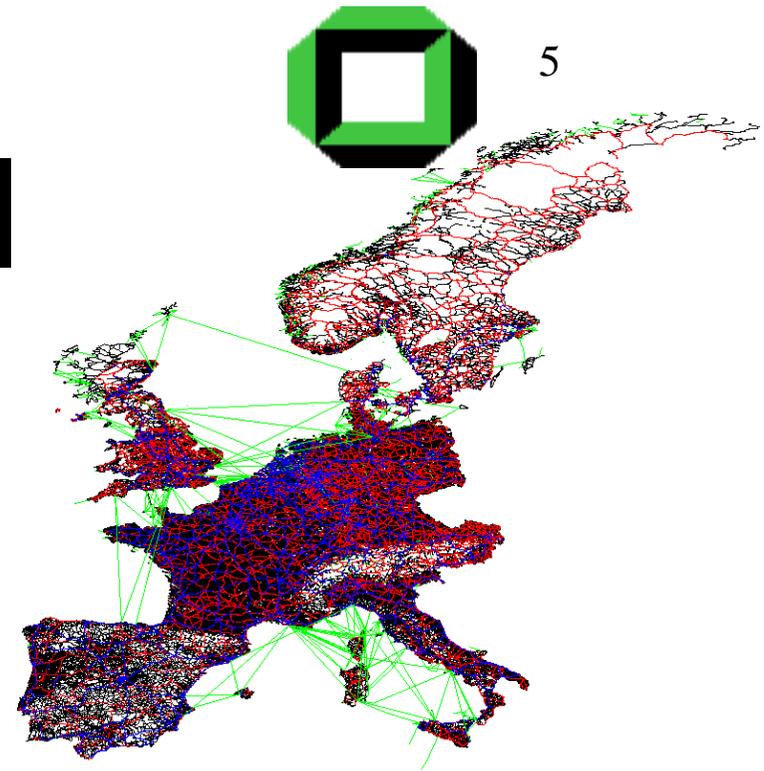


yields a **hierarchy** of highway networks

in a sense, **classify** roads / junctions by 'importance'

Highway Hierarchies

- foundation** for our other methods
- directly allows **point-to-point** queries
- 16 min** preprocessing
- 0.76 ms** to determine the path length
- 0.93 ms** to determine a complete path description
- reasonable space consumption (**68 bytes/node**)
can be reduced to **17 bytes/node**



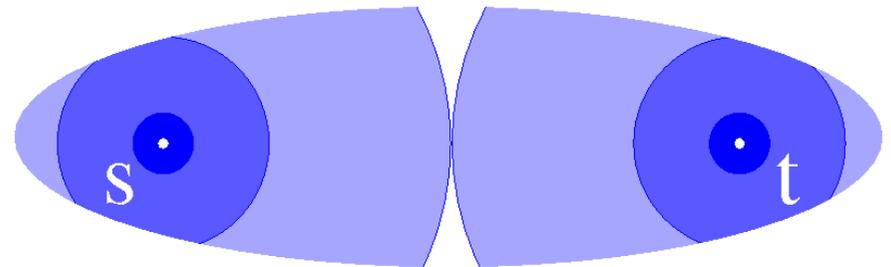


Highway Hierarchies Star

joint work with D. Delling, D. Wagner

[DIMACS Challenge 06]

- combination of highway hierarchies with **goal-directed search**
- slightly reduced query times (**0.68 ms**)
- more effective
 - for **approximate** queries or
 - when a **distance metric** instead of a travel time metric is used



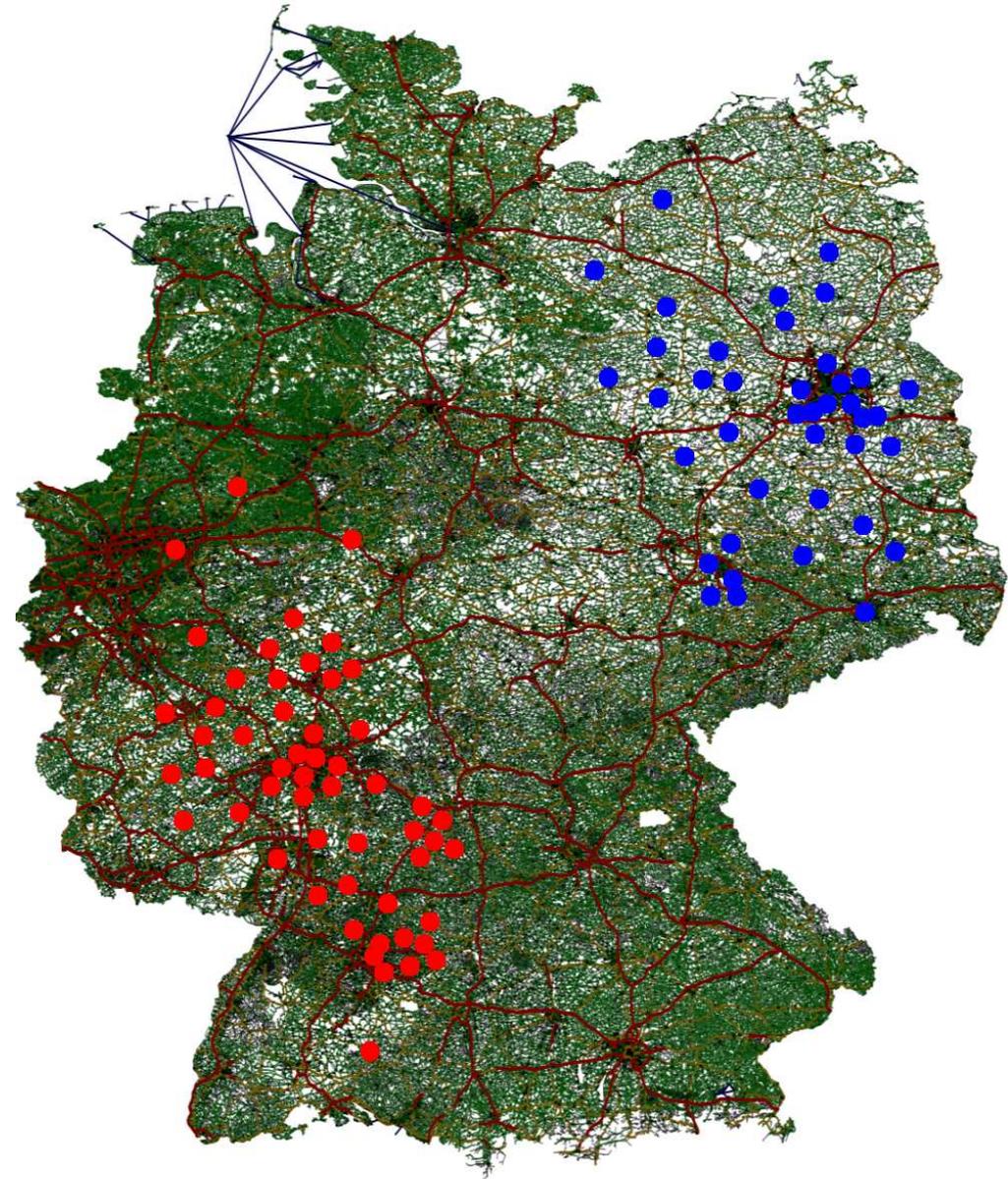
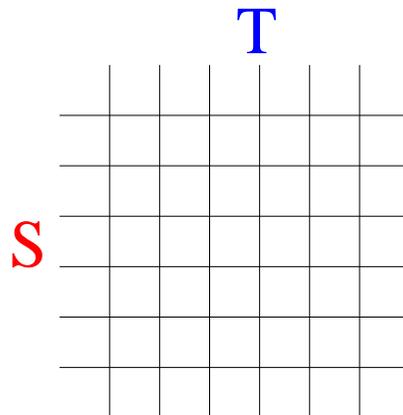


Many-to-Many Shortest Paths

joint work with S. Knopp, F. Schulz, D. Wagner

[ALENEX 07]

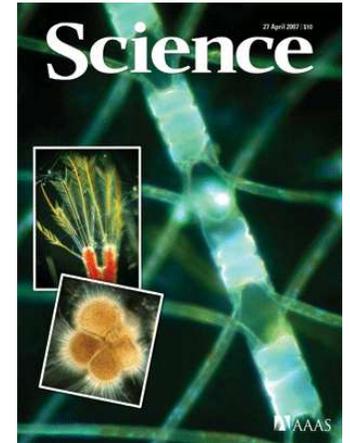
- efficient **many-to-many variant** of the highway hierarchies query algorithm
- 10 000 × 10 000 table in **one minute**





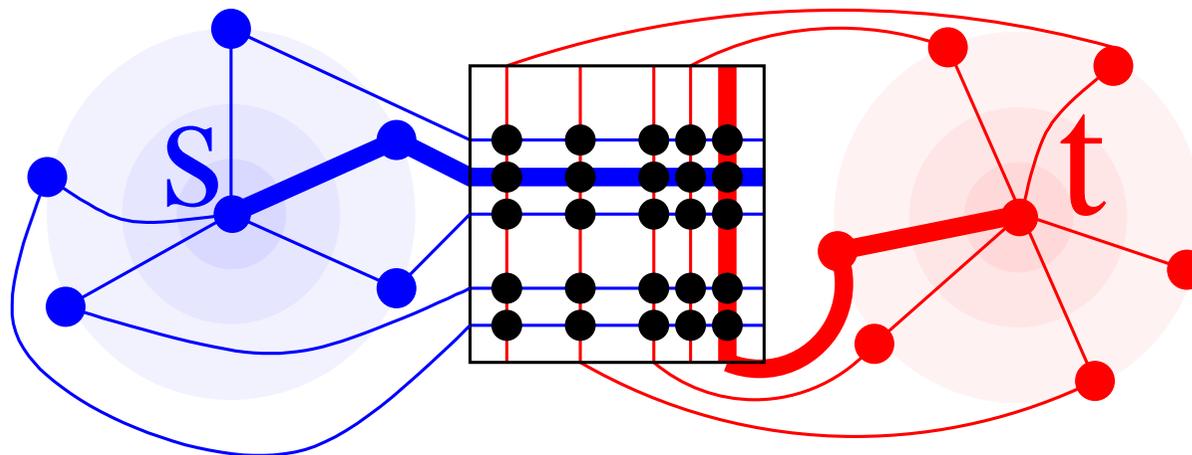
Transit-Node Routing

[DIMACS Challenge 06, ALENEX 07, Science 07]



joint work with H. Bast, S. Funke, D. Matijevic

- **very fast queries**
(down to $6 \mu s$, 1 000 000 times faster than DIJKSTRA)
- **winner** of the 9th DIMACS Implementation Challenge
- more preprocessing time (2:44 h) and space (251 bytes/node) needed

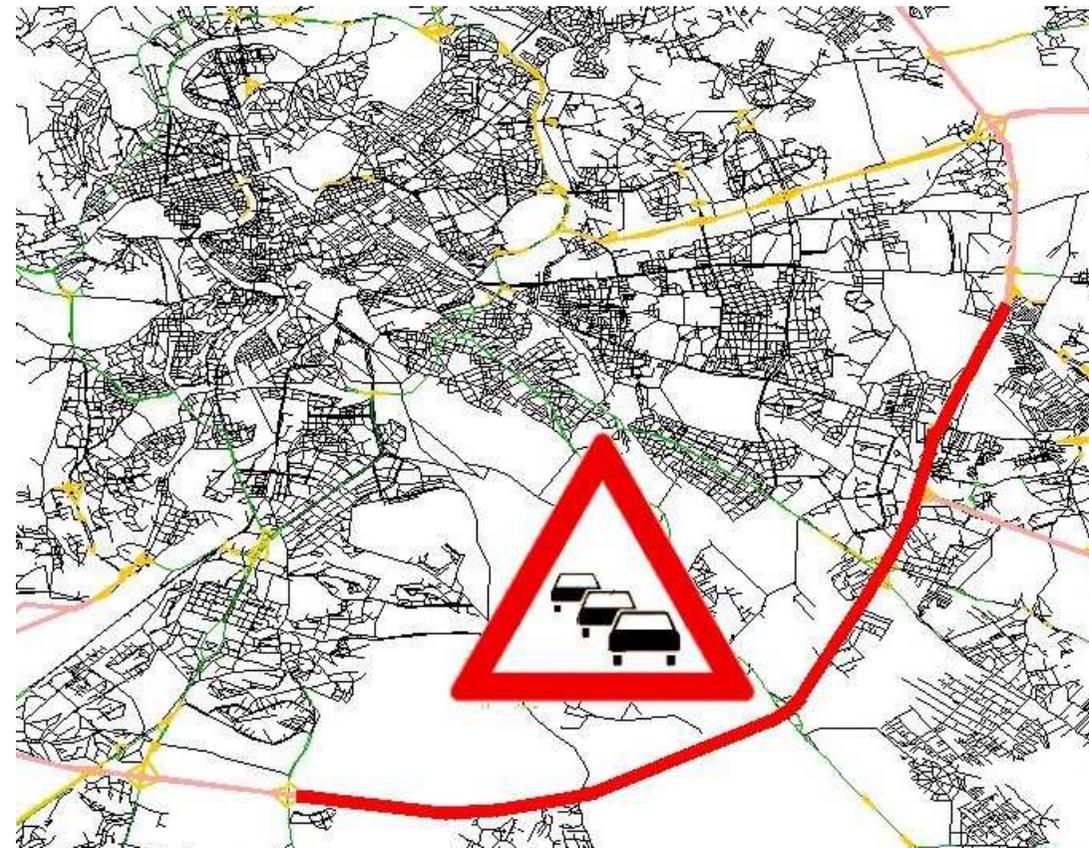




Highway-Node Routing

[WEA 07]

- performance similar to highway hierarchies
- outstandingly low** space requirements
- conceptually **very simple**
- handles **dynamic scenarios**

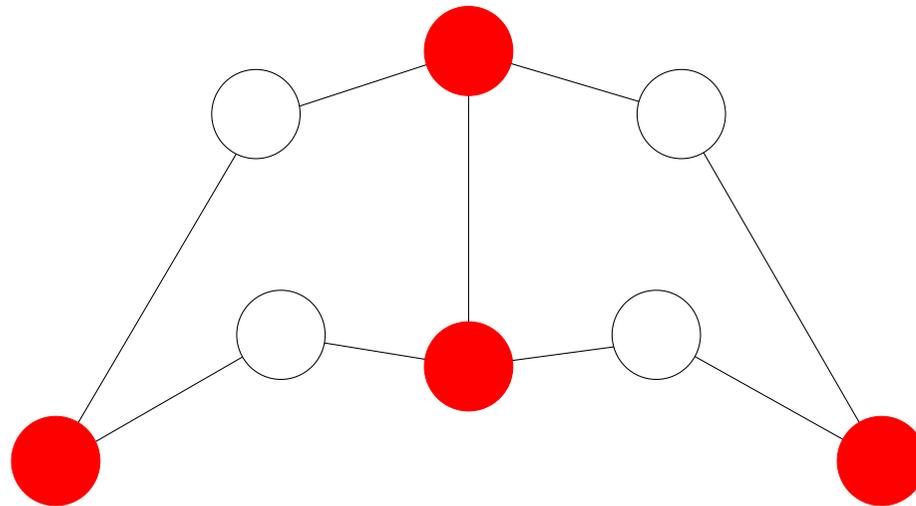




Overlay Graph

[Holzer, Schulz, Wagner, Weihe, Zaroliagis 2000–2007]

- graph $G = (V, E)$ is given
- select node subset $S \subseteq V$

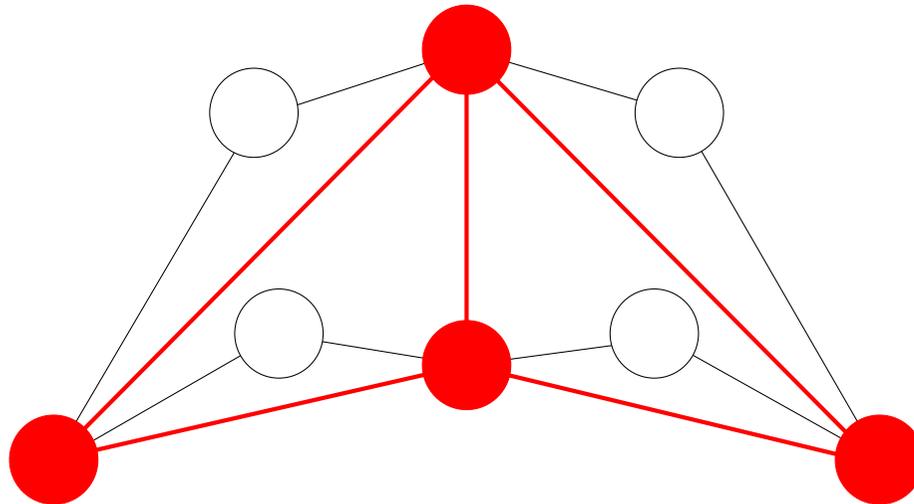




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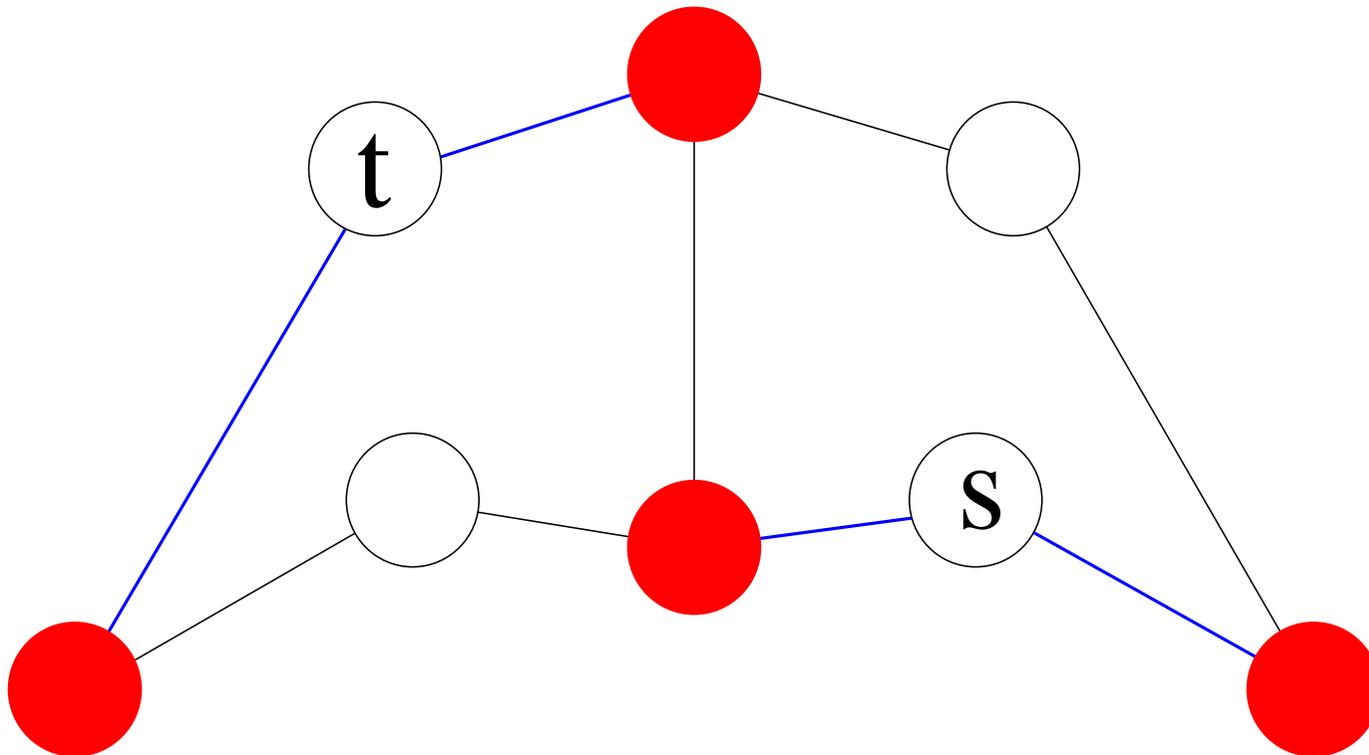
- overlay graph $G' := (S, E')$ where

$$E' := \{(s, t) \in S \times S \mid \text{no inner node of the shortest } s\text{-}t\text{-path belongs to } S\}$$



Query

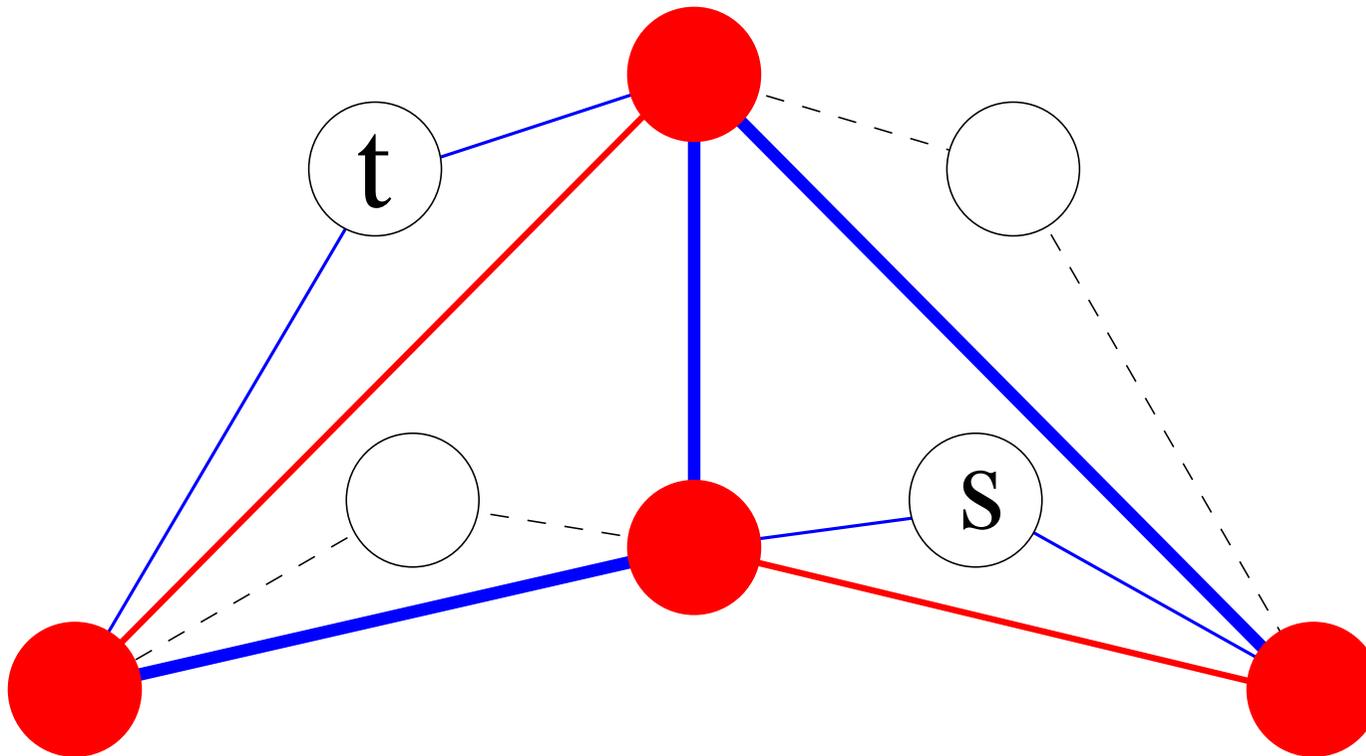
- bidirectional
- perform search in G till search trees are covered by nodes in S





Query

- bidirectional
- perform search in G till search trees are covered by nodes in S
- continue search only in G'





Static Highway-Node Routing

- extend ideas from
 - multi-level **overlay graphs** [HolzerSchulzWagnerWeiheZaroliagis00–07]
 - highway hierarchies [SS05–06]
 - transit node routing [BastFunkeMatijevicSS06–07]

- use highway hierarchies to **classify** nodes by ‘**importance**’
i.e., select node sets $S_1 \supseteq S_2 \supseteq S_3 \dots$
(crucial **distinction** from previous **separator-based** approach)

- construct **multi-level overlay graph**

- perform multi-level query



Static Highway-Node Routing

- extend ideas from
 - multi-level **overlay graphs** [HolzerSchulzWagnerWeiheZaroliagis00–07]
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 - transit node routing [BastFunkeMatijevicSS06–07]

- use highway hierarchies to **classify** nodes by ‘**importance**’
i.e., select node sets $S_1 \supseteq S_2 \supseteq S_3 \dots$ 16 min
(crucial **distinction** from previous **separator-based** approach)

- construct **multi-level overlay graph** 3 min, 8 bytes/node

- perform multi-level query 1.1 ms

(experiments with a European road network with \approx 18 million nodes)

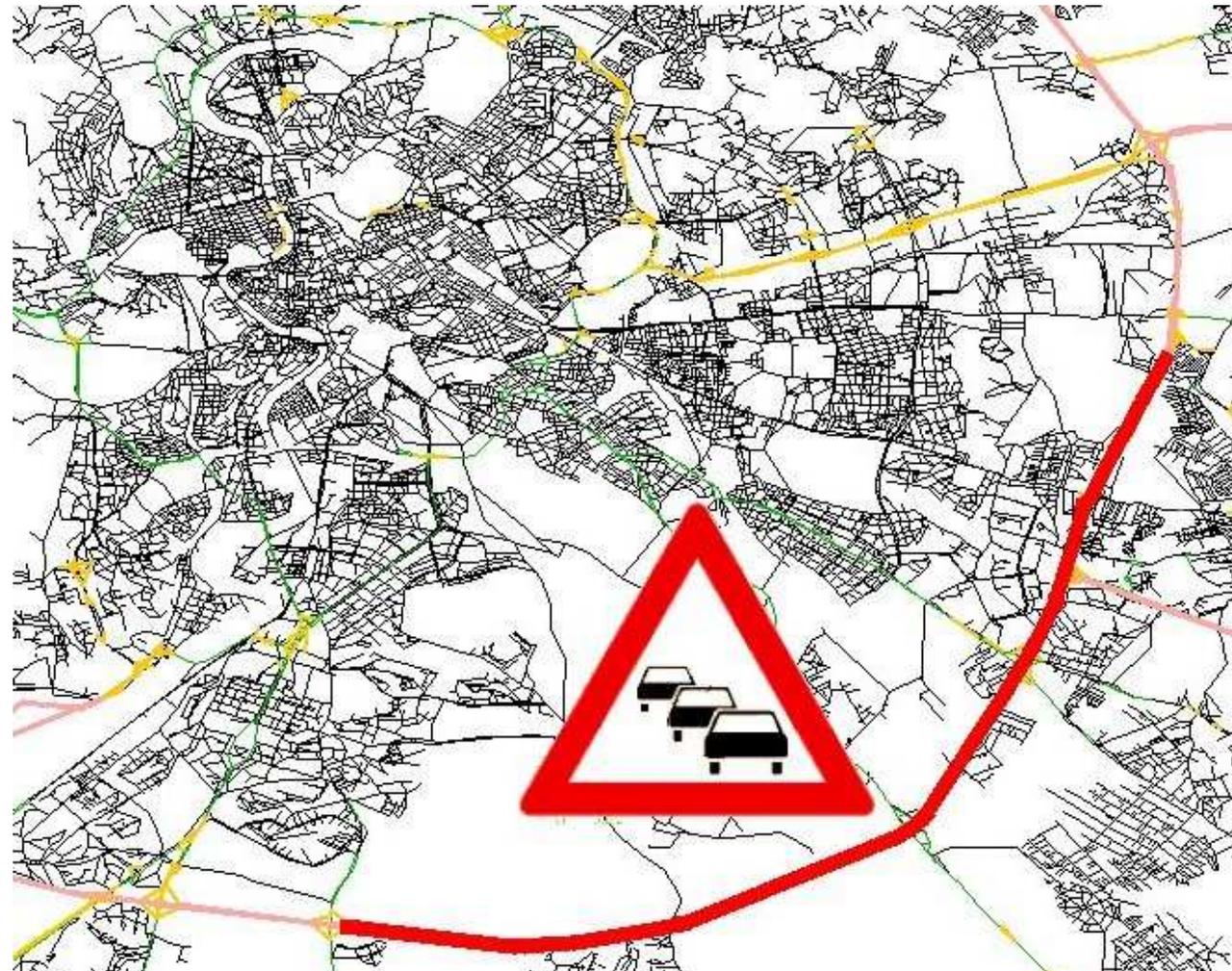


Dynamic Scenarios

- change entire **cost function**
(e.g., use different speed profile)



- change a **few edge weights**
(e.g., due to a traffic jam)





Dynamic Highway-Node Routing

change entire **cost function**



keep the node sets $S_1 \supseteq S_2 \supseteq S_3 \dots$

recompute the overlay graphs

speed profile	default	fast car	slow car	slow truck	distance
constr. [min]	1:40	1:41	1:39	1:36	3:56
query [ms]	1.17	1.20	1.28	1.50	35.62
#settled nodes	1 414	1 444	1 507	1 667	7 057



Dynamic Highway-Node Routing

change a **few edge weights**



- server scenario:** if something changes,
 - **update** the preprocessed data structures
 - answer **many** subsequent queries very **fast**

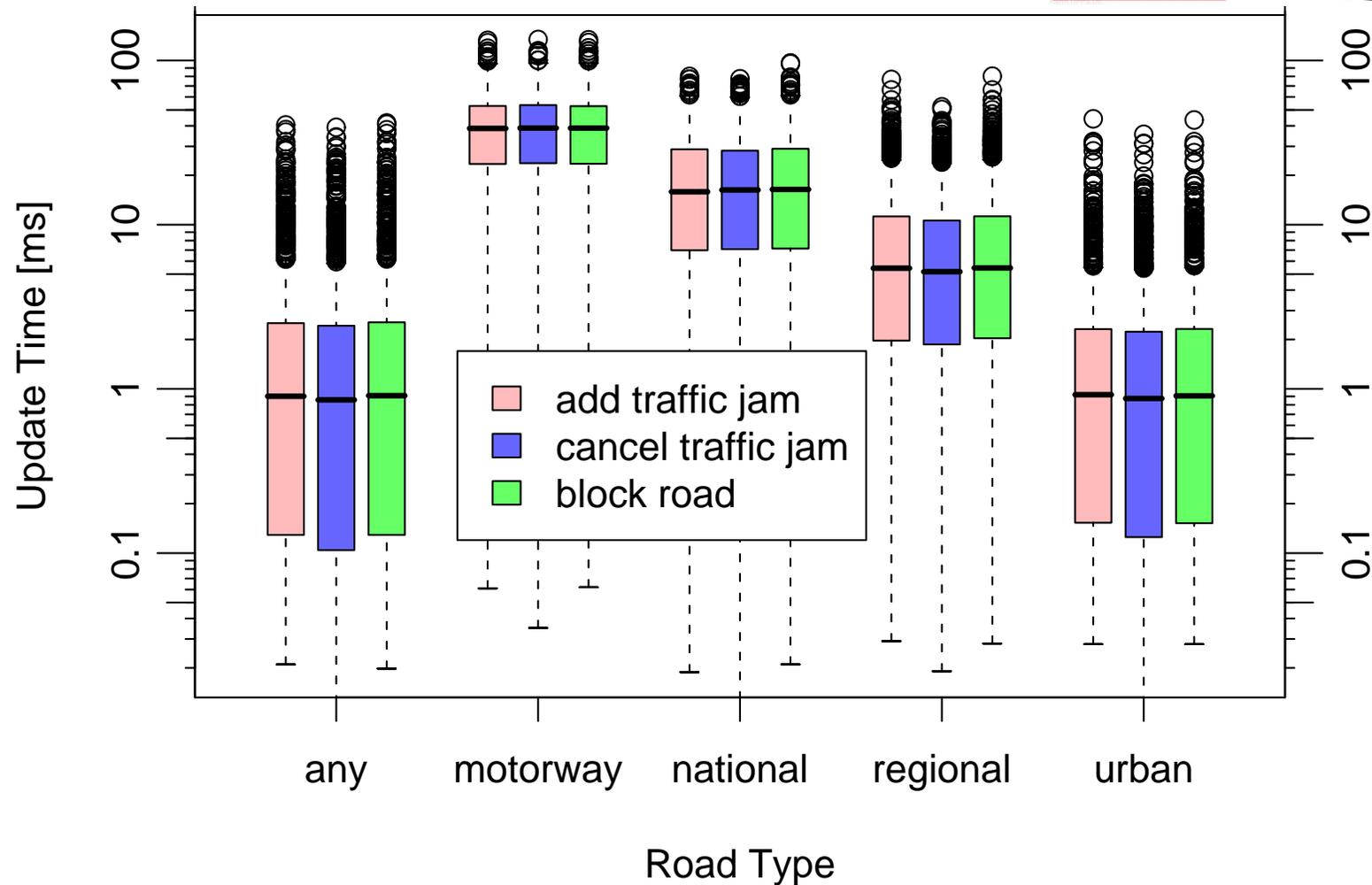
- mobile scenario:** if something changes,
 - it **does not pay** to update the data structures
 - perform **single** ‘prudent’ query that **takes changed situation into account**





Dynamic Highway-Node Routing

change a **few edge weights**, server scenario



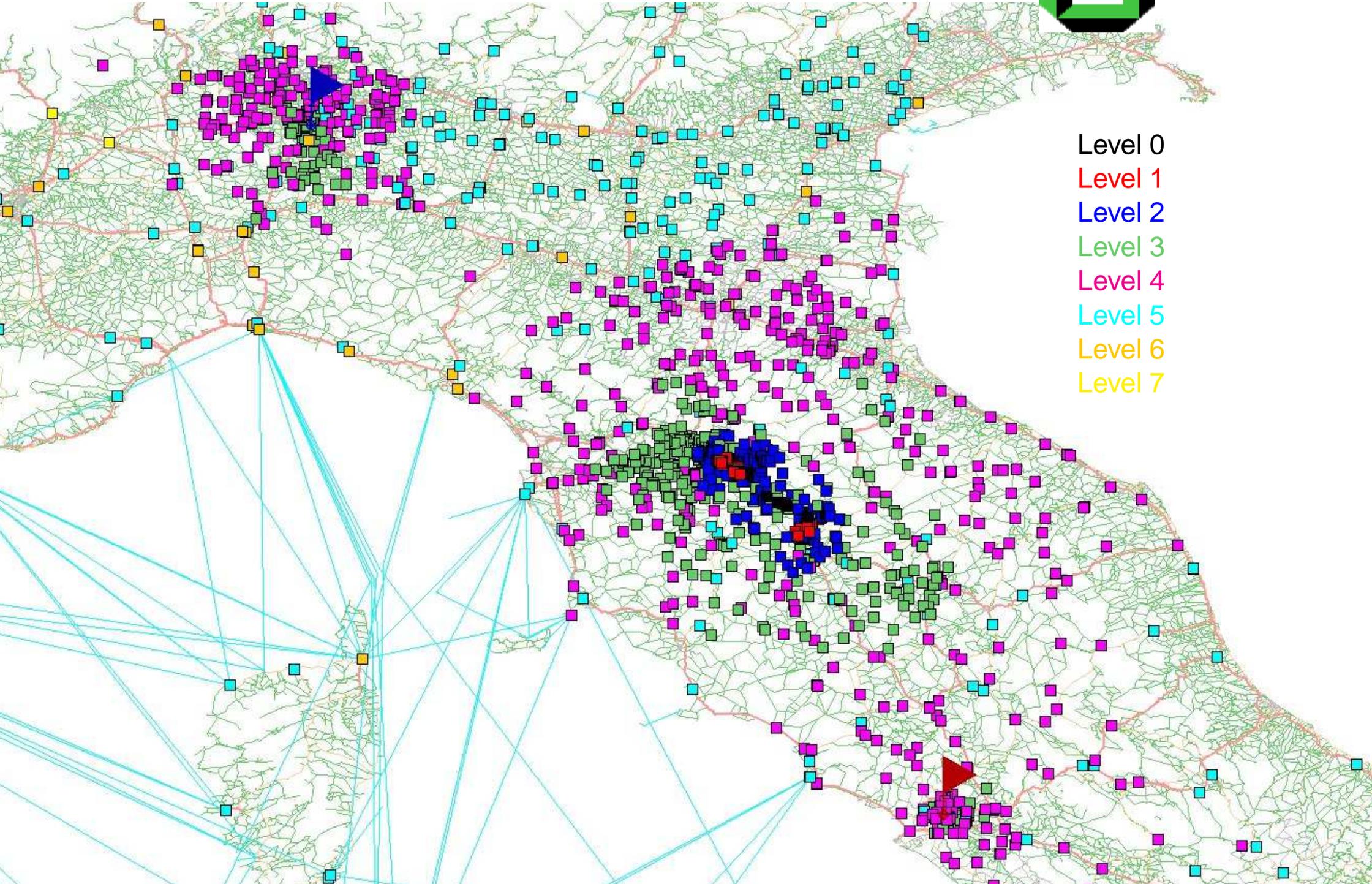


Dynamic Highway-Node Routing

change a **few edge weights**, mobile scenario



- keep** the node sets $S_1 \supseteq S_2 \supseteq S_3 \dots$
- keep** the overlay graphs
- use auxiliary data to determine for each node u a **reliable level** $r(u)$
- during a query, at node u
 - **do not use** edges that have been created in some **level** $> r(u)$
 - instead, **downgrade** the search to **level** $r(u)$





Summary

Highway Hierarchies: fast queries, fast preprocessing, low space, few tuning parameters, **basis** for many-to-many and transit-node / highway-node routing.

Many-to-Many: huge distance tables are tractable, subroutine for transit-node routing.

Transit-Node Routing: **fastest** routing so far.

Highway-Node Routing: ‘simpler’ highway hierarchies, fast queries, very low space, efficiently **dynamizable**.

